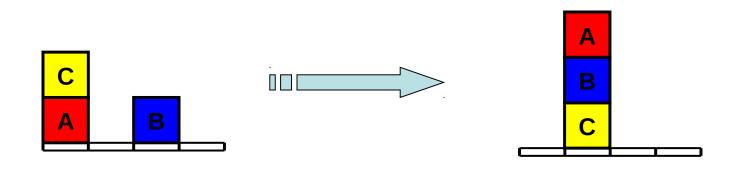
Planning Problems

Want a sequence of actions to turn a start state into a goal state

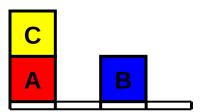


Unlike generic search, states and actions have internal structure, which allows better search methods

This slide deck courtesy of Dan Klein at UC Berkeley

State Space

On(C, A)
On(A, Table)
On(B, Table)
Clear(C)
Clear(B)



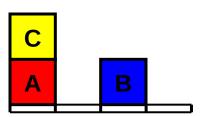
Representation

States described by propositions or ground predicates Sparse encoding (database semantics): all unstated literals are false

Unique names: each object has its own single symbol

Actions

On(C, A)
On(A, Table)
On(B, Table)
Clear(C)
Clear(B)



ACTION: Move(b,x,y)

PRECONDITIONS: On(b,x), Clear(b), Clear(y)

POSTCONDITIONS: On(b,y), Clear(x)

 $\neg On(b,x), \neg Clear(y)$

ACTION: Move(C,A,Table)

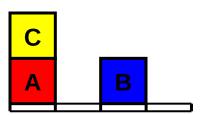
PRECONDITIONS: On(C,A), Clear(C), Clear(Table)

POSTCONDITIONS: On(C,Table), Clear(A)

 $\neg On(C,A), \neg Clear(Table)$

Actions

On(C, A)
On(A, Table)
On(B, Table)
Clear(C)
Clear(B)



ACTION: MoveToBlock(b,x,y)

PRECONDITIONS: On(b,x), Clear(b), Clear(y),

Block(b), Block(y), $(b\neq x)$, $(b\neq y)$, $(x\neq y)$

POSTCONDITIONS: On(b,y), Clear(x)

 $\neg On(b,x), \neg Clear(y)$

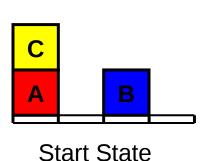
ACTION: MoveToTable(b,x)

PRECONDITIONS: On(b,x), Clear(b), Block(b), Block(x), (b \neq x)

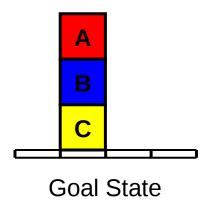
POSTCONDITIONS: On(b,Table), Clear(x)

 $\neg On(b,x)$

Start and Goal States



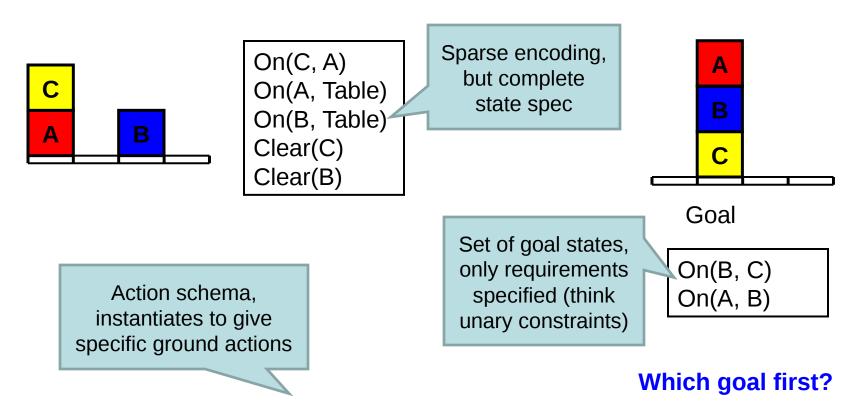
On(C, A)
On(A, Table)
On(B, Table)
Clear(C)
Clear(B)
Block(A)
...



On(B, C) On(A, B)

Important: goal satisfied by any state which entails goal list

Planning Problems



ACTION: MoveToTable(b,x)

PRECONDITIONS: On(b,x), Clear(b), Block(b), Block(x), $(b\neq x)$

POSTCONDITIONS: On(b,Table), Clear(x)

 $\neg On(b,x)$

Practice

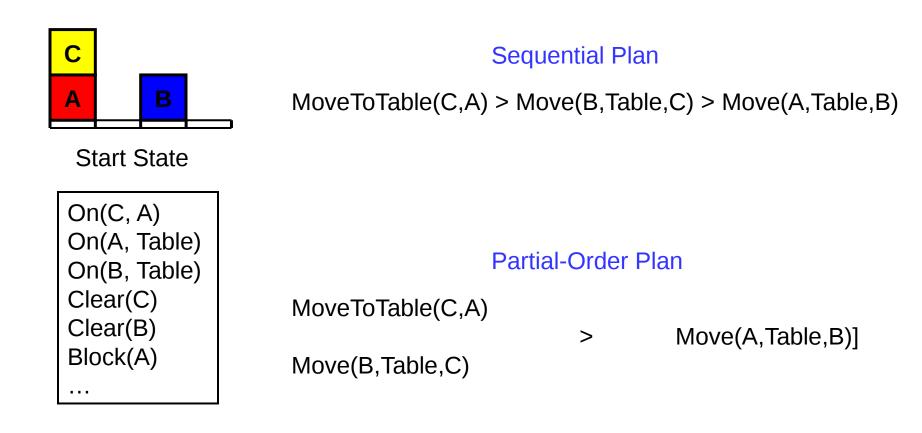
Problem 10.2: "Applicable"

Problem 10.3a,b: Representation

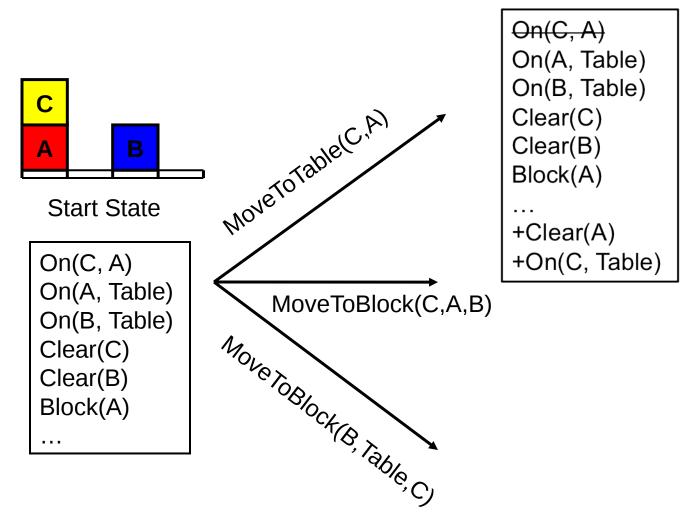
Where do they come from?

Could they be learned?

Kinds of Plans

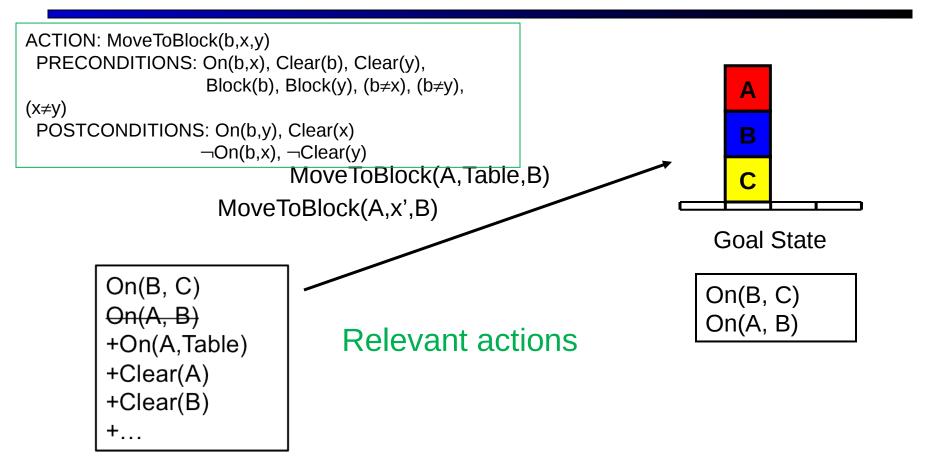


Forward Search



Applicable actions

Backward Search



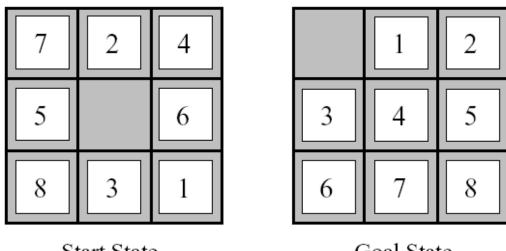
$$g' = (g - ADD(a)) \cup Precond(a)$$

Heuristics: Ignore Preconditions

Relax problem by ignoring preconditions

Can drop all or just some preconditions

Can solve in closed form or with set-cover methods



Start State

Goal State

 $Action(Slide(t, s_1, s_2),$

PRECOND: $On(t, s_1) \wedge Tile(t) \wedge Blank(s_2) \wedge Adjacent(s_1, s_2)$

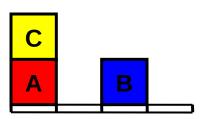
EFFECT: $On(t, s_2) \wedge Blank(s_1) \wedge \neg On(t, s_1) \wedge \neg Blank(s_2)$

Heuristics: No-Delete

Relax problem by not deleting falsified literals

Can't undo progress, so solve with hill-climbing (non-admissible)

On(C, A)
On(A, Table)
On(B, Table)
Clear(C)
Clear(B)



ACTION: MoveToBlock(b,x,y)

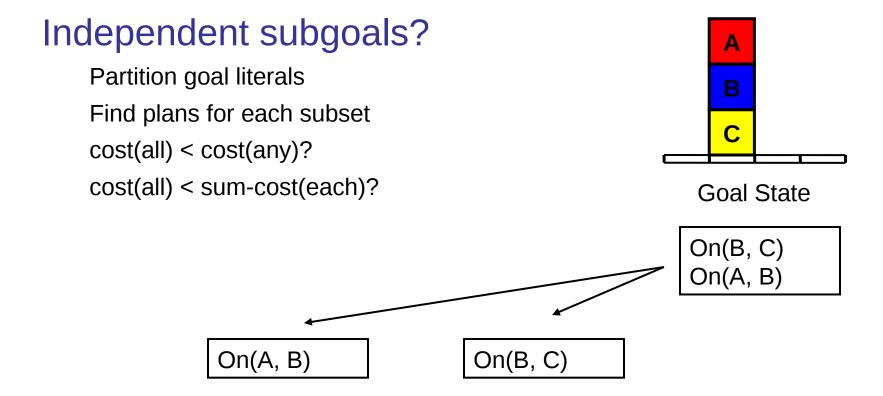
PRECONDITIONS: On(b,x), Clear(b), Clear(y),

Block(b), Block(y), $(b\neq x)$, $(b\neq y)$, $(x\neq y)$

POSTCONDITIONS: On(b,y), Clear(x)

 $\neg On(b,x), \neg Clear(y)$

Heuristics: Independent Goals



Planning "Tree"

Start: HaveCake

Goal: AteCake, HaveCake

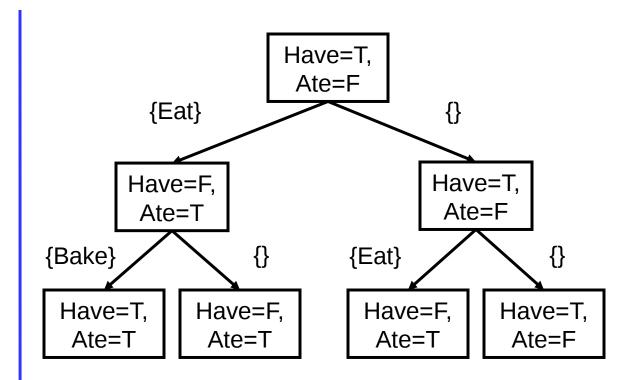
Action: Eat

> Pre: HaveCake Add: AteCake Del: HaveCake

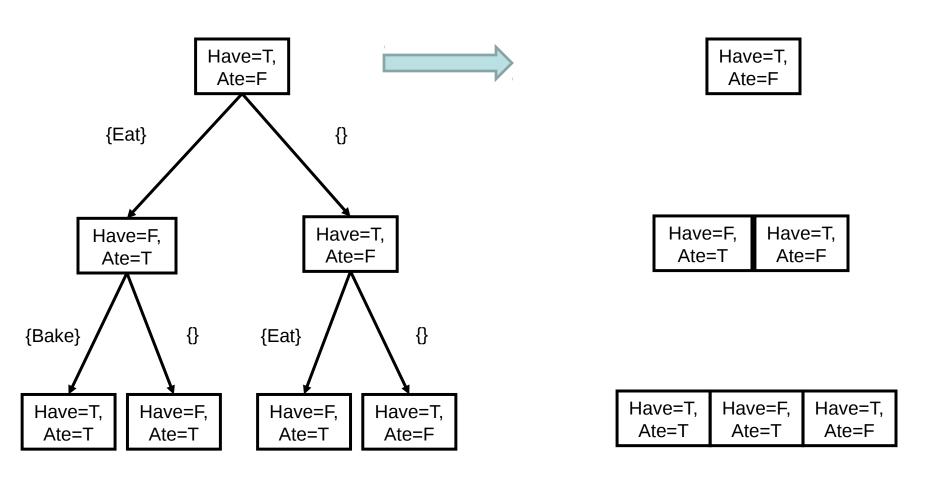
Action: Bake

Pre: ¬HaveCake

Add: HaveCake



Reachable State Sets



Approximate Reachable Sets

Have=T, Ate=F



Have={T}, Ate={F}

Have=F, Ate=T Have=T, Ate=F Have={T,F}, Ate={T,F} (Have, Ate) not (T,T) (Have, Ate) not (F,F)

Have=T, Have=F, Have=T, Ate=T Ate=F

Have={T,F}, Ate={T,F} (Have, Ate) not (F,F)

Planning Graphs

Start: HaveCake

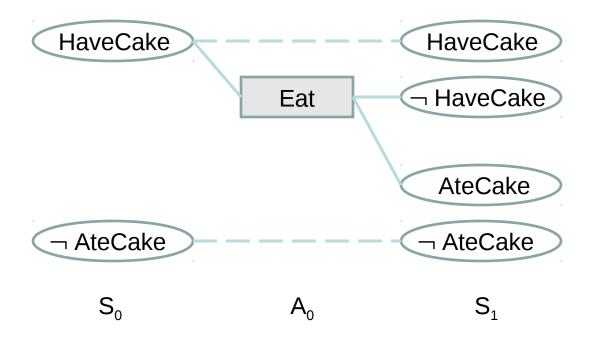
Goal: AteCake, HaveCake

Action: Eat

Pre: HaveCake Add: AteCake Del: HaveCake

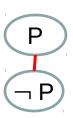
Action: Bake

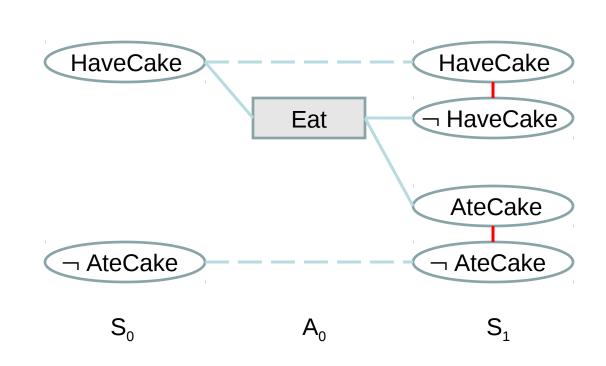
Pre: ¬HaveCake Add: HaveCake



NEGATION

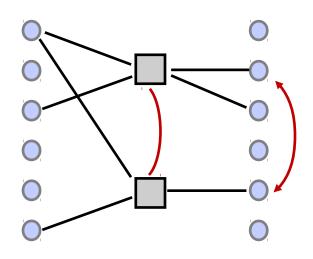
Literals and their negations can't be true at the same time

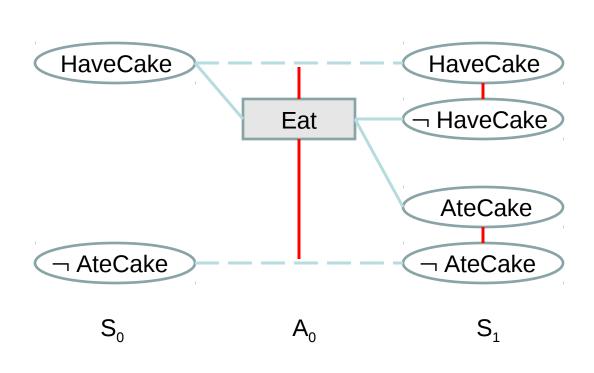




INCONSISTENT EFFECTS

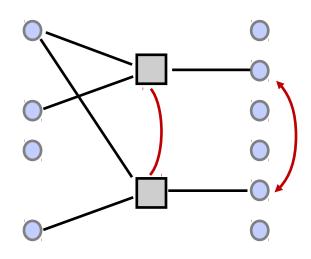
An effect of one negates the effect of the other

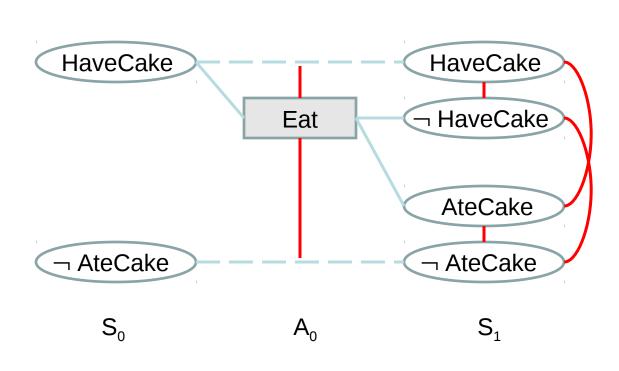




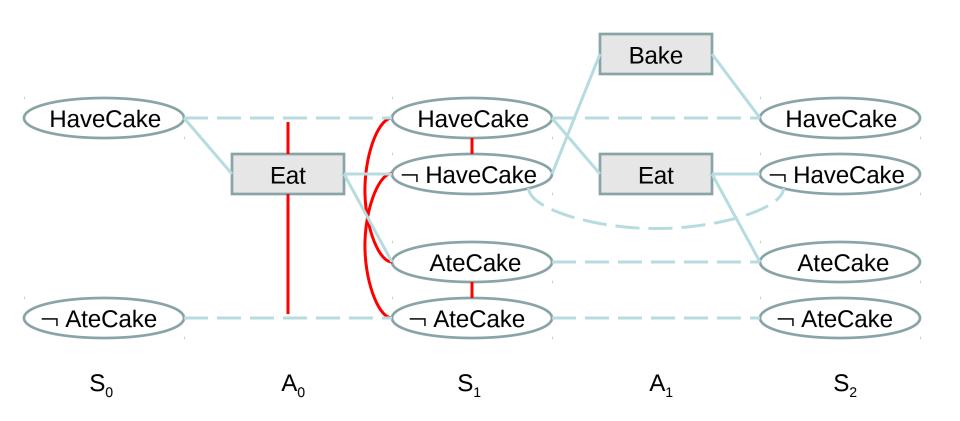
INCONSISTENT SUPPORT

All pairs of actions that achieve two literals are mutex



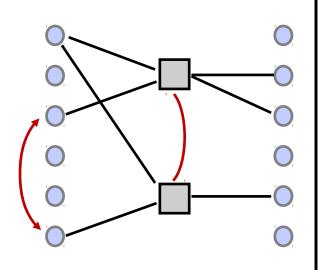


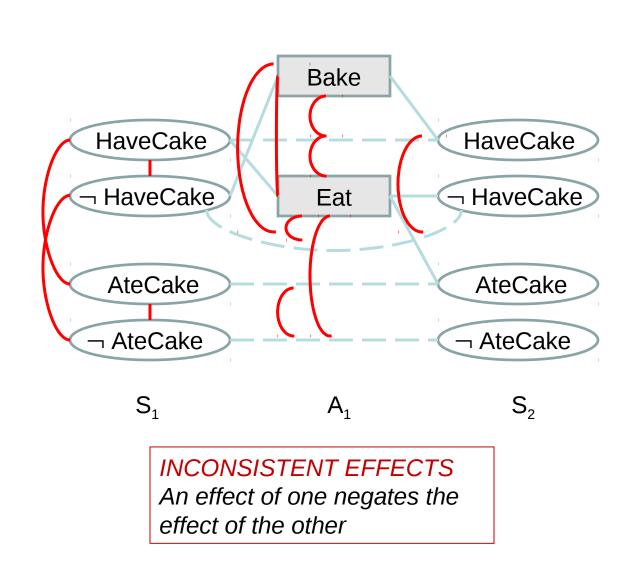
Planning Graph



COMPETITION

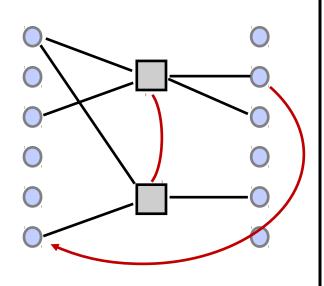
Preconditions are mutex; cannot both hold

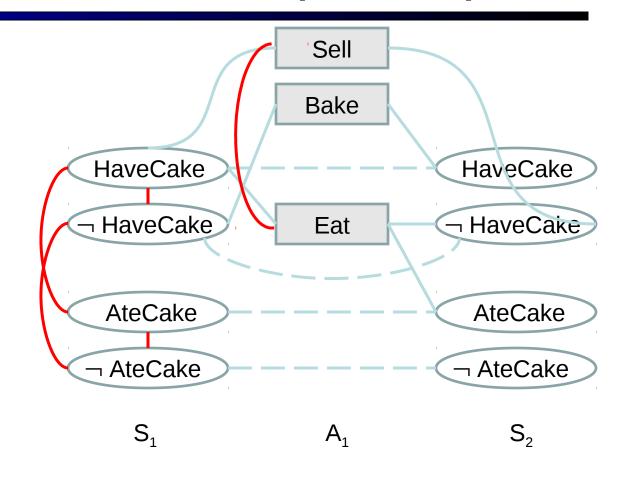




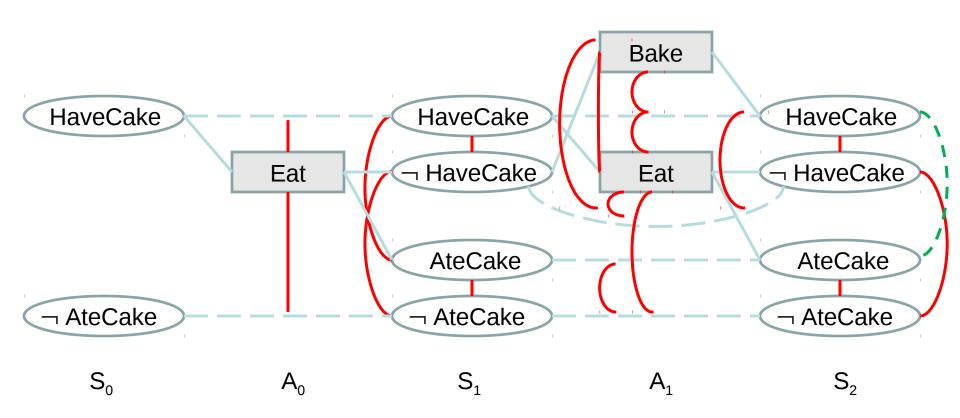
INTERFERENCE

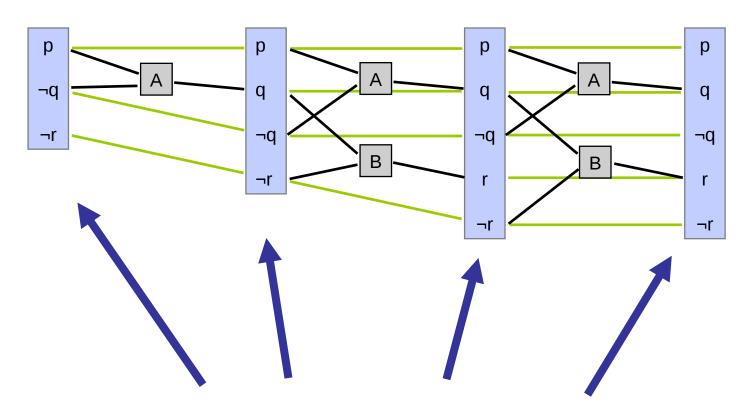
One deletes a precondition of the other



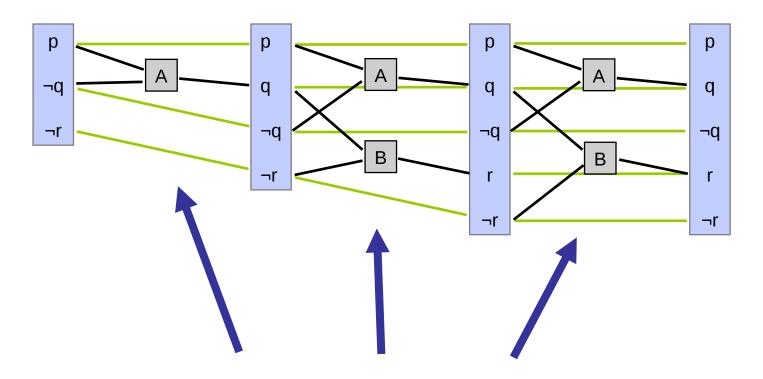


Planning Graph

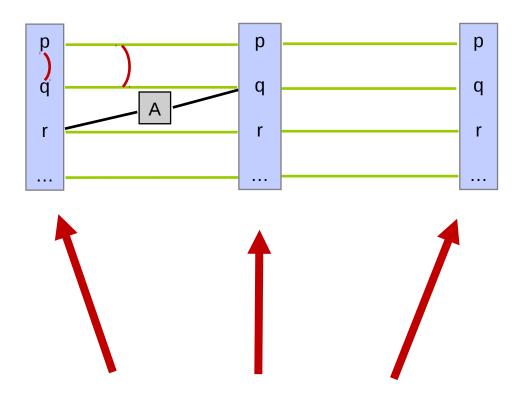




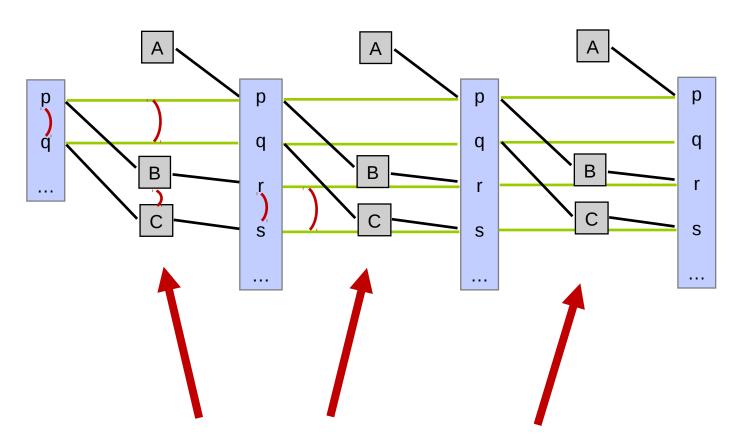
Propositions monotonically increase (always carried forward by no-ops)



Actions monotonically increase (if they applied before, they still do)



Proposition mutex relationships monotonically decrease



Action mutex relationships monotonically decrease

Claim: planning graph "levels off"

After some time k all levels are identical

Because it's a finite space, the set of literals cannot increase indefinitely, nor can the mutexes decrease indefinitely

Claim: if goal literal never appears, or goal literals never become non-mutex, no plan exists

If a plan existed, it would eventually achieve all goal literals (and remove goal mutexes – less obvious)

Converse not true: goal literals all appearing non-mutex does not imply a plan exists

Heuristics: Level Costs

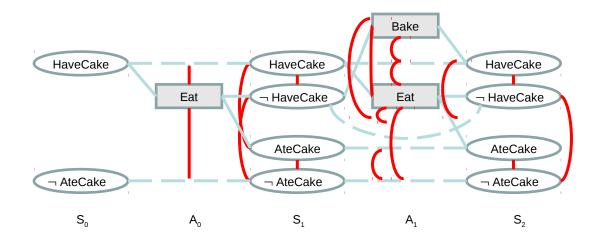
Planning graphs enable powerful heuristics

Level cost of a literal is the smallest S in which it appears

Max-level: goal cannot be realized before largest goal conjunct level cost (admissible)

Sum-level: if subgoals are independent, goal cannot be realized faster than the sum of goal conjunct level costs (not admissible)

Set-level: goal cannot be realized before all conjuncts are nonmutex (admissible)



Graphplan

Graphplan directly extracts plans from a planning graph
Graphplan searches for **layered plans** (often called parallel plans)

More general than totally-ordered plans, less general than partially-ordered plans

A layered plan is a sequence of **sets** of actions

actions in the same set must be compatible all sequential orderings of compatible actions gives same result



Layered Plan: (a two layer plan)

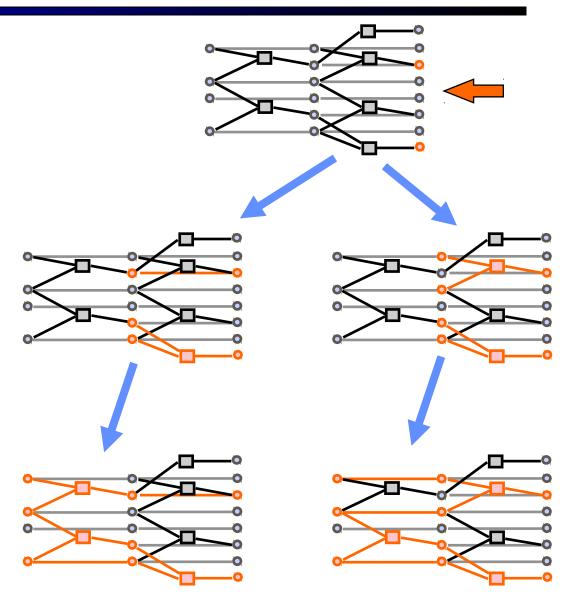
Solution Extraction: Backward Search

Search problem:

Start state: goal set at last level Actions: conflict-free ways of achieving the current goal set Terminal test: at S₀ with goal set entailed by initial planning state

Note: may need to start much deeper than the leveling-off point!

Caching, good ordering is important



Scheduling

In real planning problems, actions take time, resources

Actions have a duration (time to completion, e.g. building)

Actions can consume (or produce) resources (or both)

Resources generally limited (e.g. minerals, SCVs)

Simple case: known (partial) plan, just need to schedule

Even simpler: no resources, just ordering and duration

JOBS

[AddEngine1 < AddWheels1 < Inspect1] [AddEngine2 < AddWheels2 < Inspect2]

RESOURCES

EngineHoists (1)
WheelStations (1)
Inspectors (2)

ACTIONS

AddEngine1: DUR=30, USE=EngHoist(1)
AddEngine2: DUR=60, USE=EngHoist(1)
AddWheels1: DUR=30, USE=WStation(1)
AddWheels2: DUR=15, USE=WStation(1)
Inspect1: DUR=10, USE=Inspectors(1)
Inspect2: DUR=10, USE=Inspectors(1)

Resource-Free Scheduling

JOBS

[AddEngine1 < AddWheels1 < Inspect1] [AddEngine2 < AddWheels2 < Inspect2]

RESOURCES

EngineHoists (1) WheelStations (1) Inspectors (2)

ACTIONS

AddEngine1: DUR=30, USE=EngHoist(1)
AddEngine2: DUR=60, USE=EngHoist(1)
AddWheels1: DUR=30, USE=WStation(1)
AddWheels2: DUR=15, USE=WStation(1)
Inspect1: DUR=10, USE=Inspectors(1)
Inspect2: DUR=10, USE=Inspectors(1)

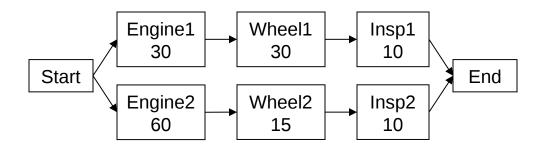
How to minimize total time?

Easy: schedule an action as soon as its parents are completed

$$ES(START) = 0$$

$$ES(a) = \max_{b:b \prec a} ES(b) + DUR(b)$$

Result:



Resource-Free Scheduling

JOBS

[AddEngine1 < AddWheels1 < Inspect1] [AddEngine2 < AddWheels2 < Inspect2]

RESOURCES

EngineHoists (1) WheelStations (1) Inspectors (2)

ACTIONS

AddEngine1: DUR=30, USE=EngHoist(1)
AddEngine2: DUR=60, USE=EngHoist(1)
AddWheels1: DUR=30, USE=WStation(1)
AddWheels2: DUR=15, USE=WStation(1)
Inspect1: DUR=10, USE=Inspectors(1)
Inspect2 DUR=10, USE=Inspectors(1)

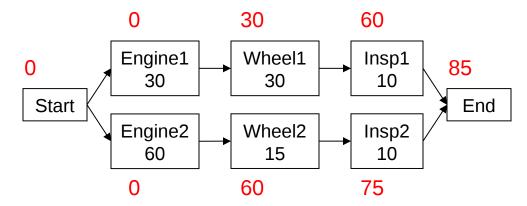
Note there is always a critical path All other actions have slack

Can compute slack by computing latest start times

$$LS(END) = ES(END)$$

$$LS(a) = \min_{b: a \prec b} LS(b) - DUR(a)$$

Result:



Adding Resources

For now: consider only released (non-consumed) resources

View start times as variables in a CSP

Before: conjunctive linear constraints

$$\forall b: b \prec a \quad ES(a) \geq ES(b) + DUR(b)$$

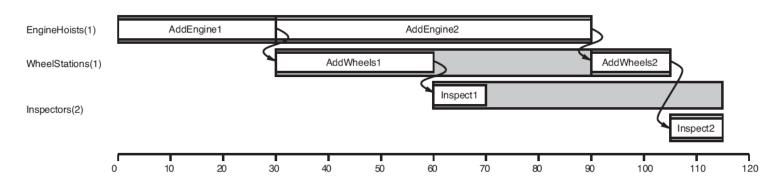
Now: disjunctive constraints (competition)

if competing (a, b)

$$ES(a) \ge ES(b) + DUR(b) \lor$$

$$ES(b) \ge ES(a) + DUR(a)$$

In general, no efficient method for solving optimally



Adding Resources

One greedy approach: min slack algorithm

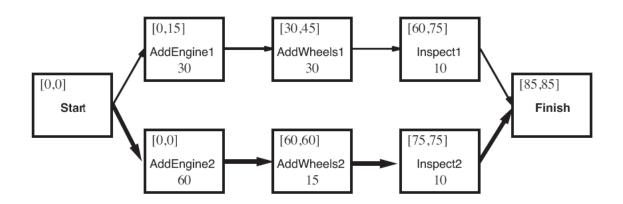
Compute ES, LS windows for all actions

Consider actions which have all preconditions scheduled

Pick the one with least slack

Schedule it as early as possible

Update ES, LS windows (recurrences now must avoid reservations)



Resource Management

Complications:

Some actions need to happen at certain times

Consumption and production of resources

Planning and scheduling generally interact